Liquidity decisions and the timing of payments Rafael Jiménez September 2015



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Contents

1 What? Explain behavior in an interbank payment system

Why? Crucial to understanding risks

3 How? The model

4 Lessons? Results and next steps

Figure : SPEI (Mex. RTGS) network



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Output: Simulator of a Real-Time Gross Settlement (RTGS) system Timing of payments

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Explain behavior of participants in an interbank payment system

Behavior: liquidity decisions, sent payments

Natural framework: Game. Interconnectedness!

Output: Simulator of a Real-Time Gross Settlement (RTGS) system Timing of payments Intraday liquidity demand

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To ensure the sound functioning of payment systems $\rightarrow \textbf{adequate}$ risk-assessment



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Participants' decision-making process: crucial to understanding risks



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Participants' decision-making process: crucial to understanding risks

'Fixed' behavior may not be an adequate assumption:

 $\operatorname{action}_{t}^{i}(\operatorname{history}_{t}) \neq \operatorname{action}_{t}^{i}(\operatorname{history}_{t}')$

Measuring liquidity demand and the timing of payments

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From a central bank's perspective, interesting to measure (and explain) the demand for intraday liquidity and the timing of payments

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Overview

Extensive game (over each settlement cycle) with imperfect information (future incoming payments and payment requests)

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Pioneering work: competitive multi-agent model of interbank payment systems (Galbiati & Soramäki, 2011).

Model description (1/4)

Set of participants $\mathcal{N} = \{1, 2, \dots, N\}$

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Participants face budget constraint. Funding sources: Intraday liquidity provided by the central bank L_t^i Positive account balances they save from previous period B_t^i Received payments from others on previous period R_t^i

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How? Model description (2/4)

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Objective: fulfill (exogenous) payment requests timely, with minimum liquidity costs. Unfulfilled payment requests at period *t*: $\mathcal{O}_t^i = \left(\mathcal{O}_{t,k}^{ij}\right)_k$

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Reputation costs of period *t*:

$$\mathsf{Reputation}_{t}^{i} = \sum_{k} \kappa O_{t,k}^{ij} \left(1 - x_{k}^{i} \right)$$

where $x_k^i = 1$ if payment is settled, $x_k^i = 0$ otherwise, κ is the "interest cost" of delaying payments (Galbiati & Soramäki, 2011).

Model description (3/4)

Settled payments (first-in-first-out algorithm):

$$x_k^i = 1 \iff \sum_{p=1}^k O_{t,k}^{ij} \le B_t^i + L_t^i + R_t^i$$

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$$Liquidity'_t = qL'_t + F\mathbb{1}_{\left\{L'_t > 0\right\}}$$

Return function for $t \in \{0, \dots, T-1\}$:

$$\mathsf{Return}_t^i\left(L_t^i, B_t^i, R_t^i, \mathcal{O}_t^i\right) = \sum_k \kappa \mathcal{O}_{t,k}^{ij}\left(1 - x_k^i\right) + qL_t^i + F\mathbb{1}_{\left\{L_t^i > 0\right\}}$$

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Terminal return function is the interest payed from missing funds (negative if surplus funds):

$$\text{Terminal}^{i}\left(\sum_{t}L_{t}^{i}, B_{T+1}^{i}, R_{T+1}^{i}\right) = r\left(\sum_{t}L_{t}^{i} - B_{T+1}^{i} - R_{T+1}^{i}\right)$$

Recursive problem Given $S_0 = (L_0^a, B_0, R_0, \mathcal{O}_0)$ and assuming individuals take as given the Markovian process R, the recursive formulation:

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Terminal condition:

$$v_{T}(S_{T}) = \beta r \left(L_{T}^{a} + \sum_{k} O_{T,k} - B_{T} - R_{T} - E_{T}(R_{T+1}) \right)$$

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Bellman-equation for t < T:

$$\begin{aligned} v_t(S_t) &= \min_{L_t} \left\{ \text{return} \left(S_t, L_t \right) + \beta \mathbf{E}_t \left(S_{t+1} \right) \right\} \\ \text{s.t.} \quad B_{t+1} &= B_t + R_t + L_t - \sum_k O_{t,k} x_{t,k} \ge 0 \\ L_{t+1}^a &= L_t^a + L_t \\ 0 \le L_t \le \overline{L} - L_t^a \\ \mathcal{O}_{t+1} &= \left(O_{t,k}, \mathcal{O}_{t+1}' \right), \ k : \sum_{p=1}^k O_{t,k} > B_t + L_t + R_t \end{aligned}$$

Equilibrium and the curse of dimensionality

If we could solve by backwards induction, we would obtain policy functions: $L_t^i(S_t^i)$, so that $S_0^i \to L_0^i(S_0^i) \to S_1^i \dots$

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Approximate dynamic programming (ADP) algorithms allow us to perform forward induction (avoid looping through all possible states)

Substitute computational problem for statistical problem: estimate value functions \overline{v}_{t+1}

Approximate dynamic programming

Estimate value function using basis functions (linear):

$$E_t(\overline{\mathbf{v}}_{t+1}) = \beta_t + \beta_1 L_t^a + \beta_2 B_t + \beta_3 R_t + \beta_4 \sum_k (\mathcal{O}_{t,k})$$

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General idea:

fix initial parameters in outer loop

sample different states (m) and calculate m parameters in inner loop

update new parameters using recursive least squares

use new parameters in outer loop

Testing the model

Two tests:

- 1. Compare the model to observed data. Choose a day at random (January 16th, 2014), and compare two measures (liquidity and timing-skewness)
- 2. Compare the model to a simulator, under a stress test. Assume that the biggest participant of a given day fails at t = 0

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Data:

89 participants

Start at 5:00 am, finish at 5:59:30 pm (1,560 30-second periods) Adapt the model to Mexico: two sources of central bank liquidity

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Time-skewness: measures concentration of payments at the end of the day (0: all payments sent during first period, 1: all payments sent in last period): $\sum_{t} (1 - \text{cum. proportion}) / (T - 1)$

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Lessons? Model vs Reality (1/2)

The model predicts reasonable time-skewness

Figure : Time-skewness by participant



Lessons? Model vs Reality (2/2)

The model overestimates liquidity (but with some correlation)



Figure : Liquidity demanded by participant

Lessons? Model vs Simulator (1/2)

The model predicts lower change in skewness

Figure : Skewness(failure)-Skewness(original) by participant



Lessons? Model vs Simulator (2/2)

The model predicts lower change in liquidity

Figure : Liquidity(failure)-Liquidity(original) by participant



Lessons?

Summary

Comparison with reality (typical day):

the model predicts reasonable time-skewness

liquidity is overestimated (but with correlation = 0.77)

	Model	Observed
Time-skewness	0.55967	0.55787
Liquidity	$1.98 imes10^{14}$	$1.23 imes 10^{11}$

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Comparison with simulator (failure of biggest participant):

the model predicts almost no change in time-skewness or liquidity simulator: delay in payments and increase in liquidity needs

	Model	Simulator
Δ Time-skewness	0.00868	0.09423
∆% Liquidity	0.04%	163.95%



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Simulate multiple days

References

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